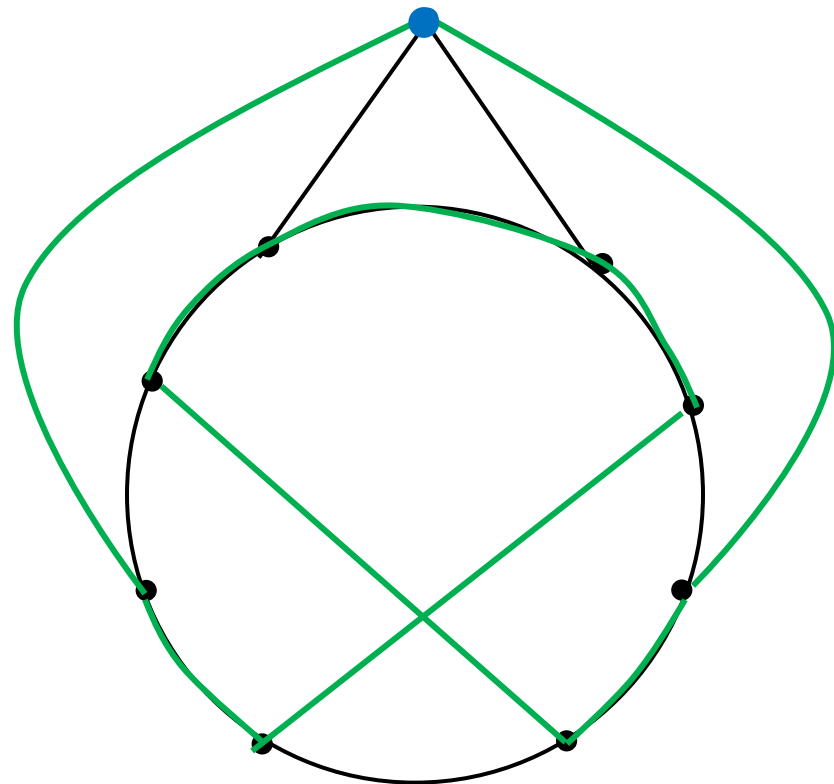
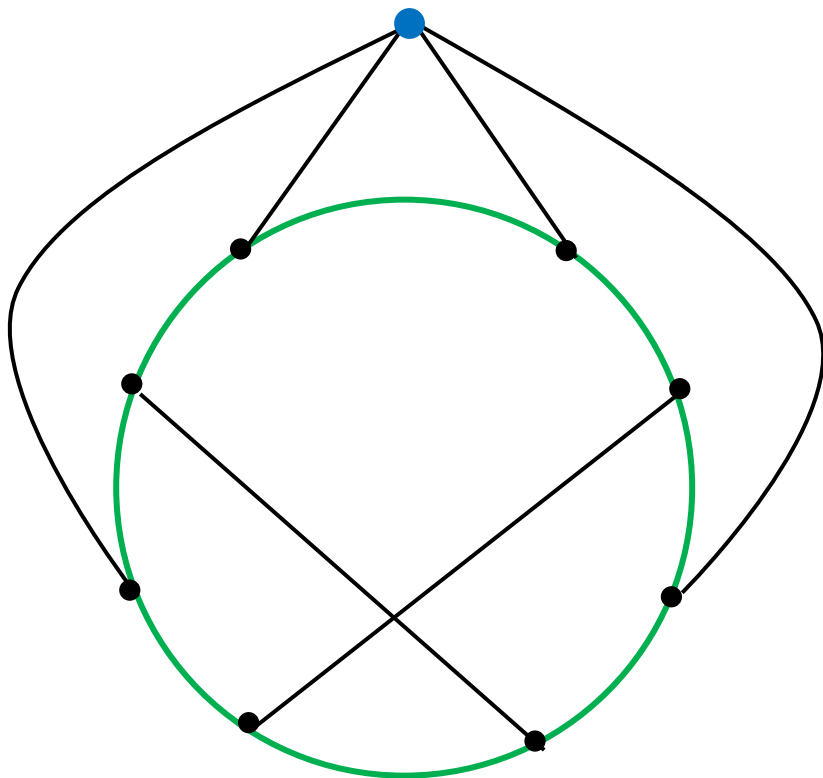


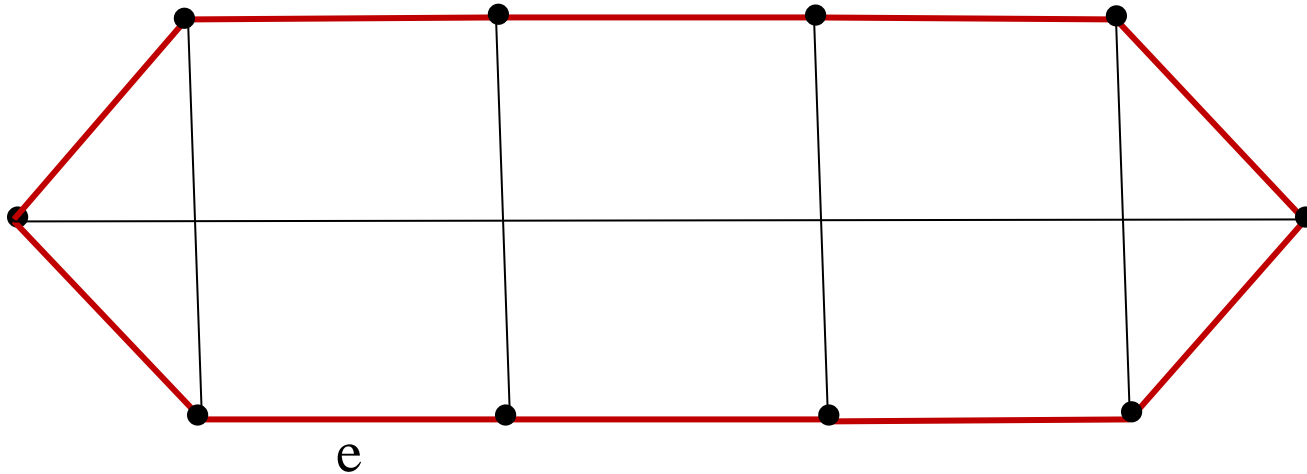
Parity Theorems about Cycles and Trees

Kathie Cameron

Wilfrid Laurier University
Waterloo, Canada



A **Hamiltonian cycle** in a graph G is a cycle containing each vertex of G



Smith's Theorem (Tutte 1946)

Let G be a graph where the degree, $d(v)$, of every vertex v is 3.

Let e be an edge of G .

Then the number of Hamiltonian cycles **containing edge e** is even.

~~Smith's Theorem~~ (Tutte 1946) Theorem (Andrew Thomason 1978)

Let G be a graph where the degree, $d(v)$, of every vertex v is ~~3~~ **odd**.

Let e be an edge of G .

Then the number of Hamiltonian cycles containing edge e is even.

Exchange Graphs

The idea:

For any graph, the number of vertices of odd degree is even.

To prove that the number of **desired structures** is even,
construct a graph **X** such that

desired structures \leftrightarrow odd-degree vertices of **X**

Then, given one **desired structure**, to find another **desired structure**,
walk in the **exchange graph X** from the given odd-degree vertex to
another odd-degree vertex

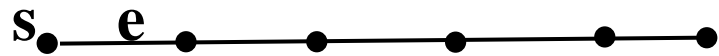
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Let G be a graph where the degree, $d(v)$, of every vertex v is odd.

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Proof.



Hamiltonian path starting with e

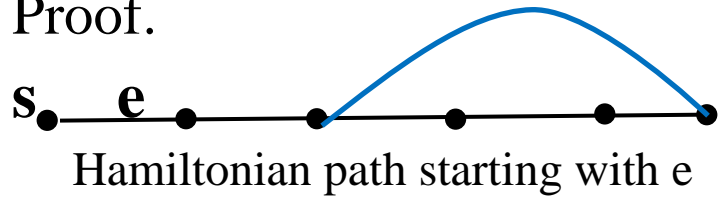
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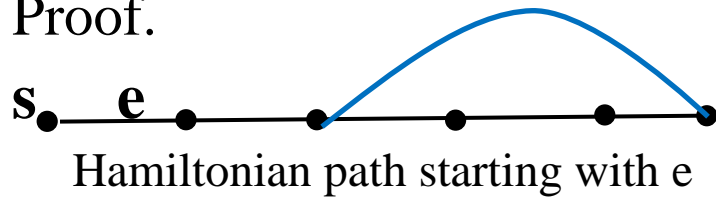
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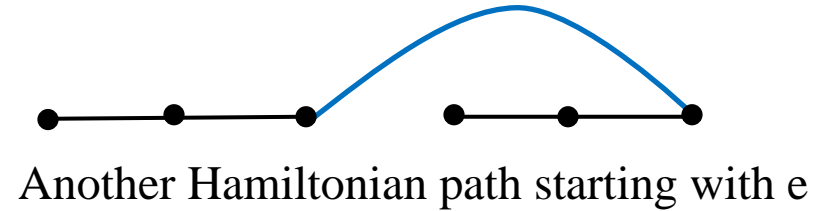
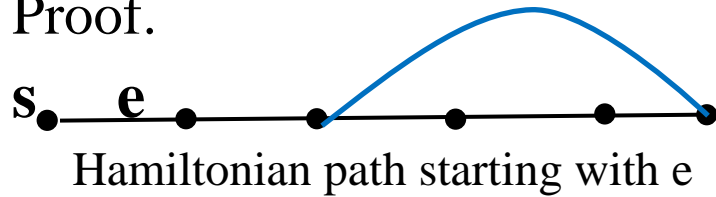
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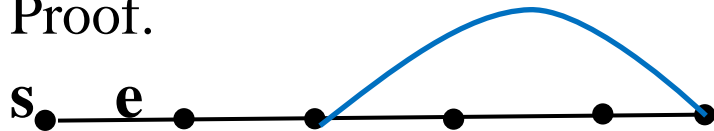
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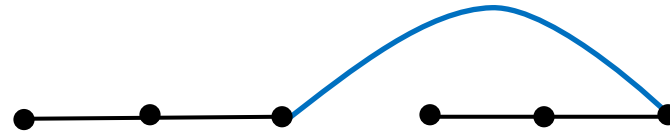
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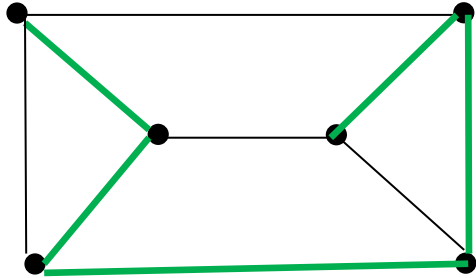
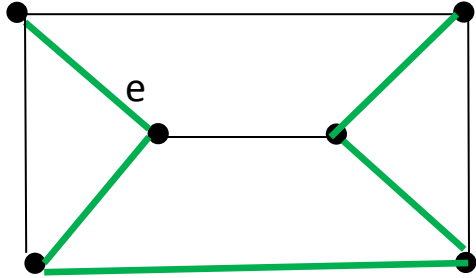


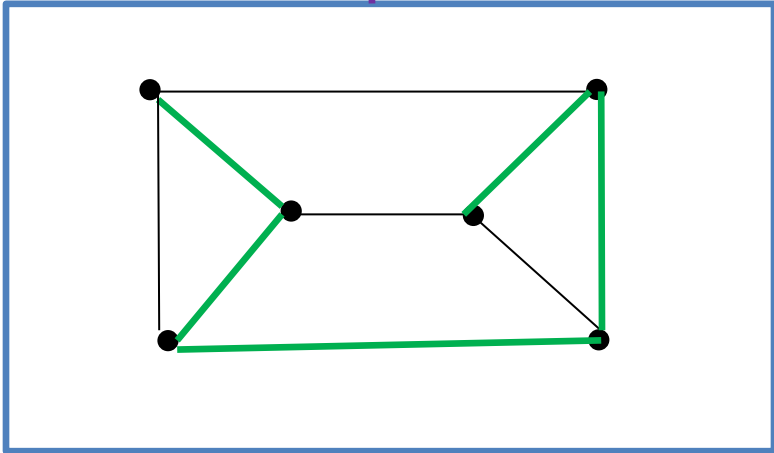
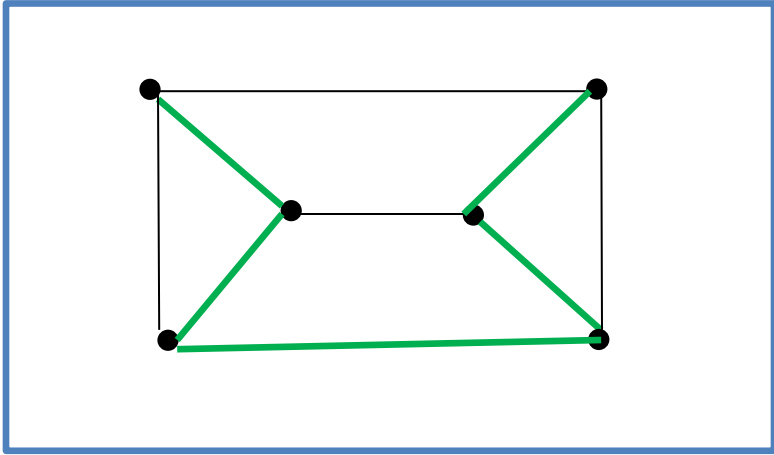
Another Hamiltonian path starting with e

Create the **exchange graph $X(G)$** :

Vertices are: Hamiltonian paths starting with e

Join two vertices of $X(G)$ if the Hamiltonian paths they correspond to are related by an exchange as above





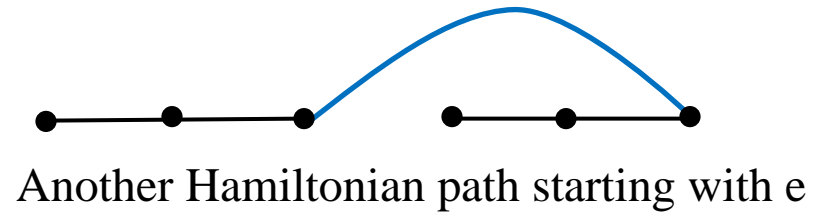
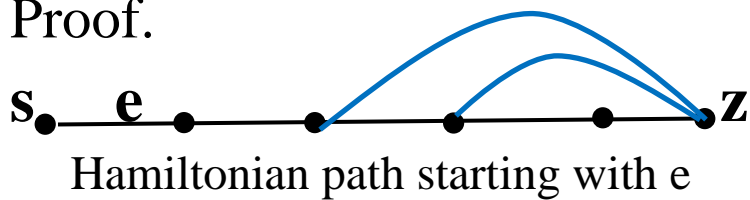
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Degrees in $X(G)$:

Let P be a vertex of $X(G)$, i.e., a Hamiltonian path starting with e and ending at z

$\text{degree}_{X(G)}(P) =$
if P is not extendible to a Hamiltonian cycle
if P is extendible to a Hamiltonian cycle

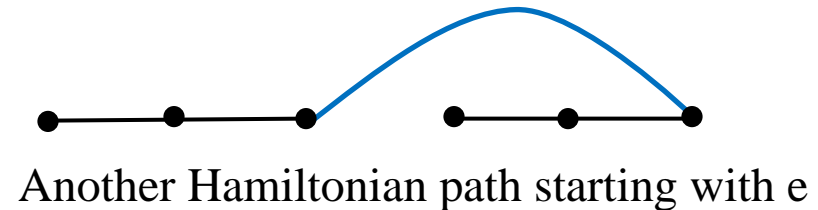
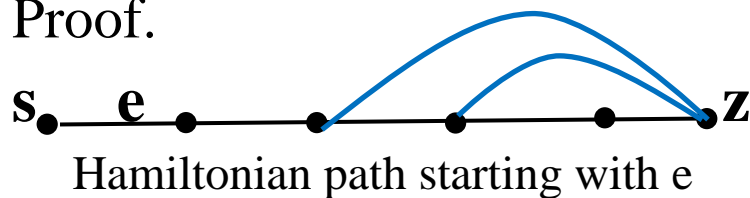
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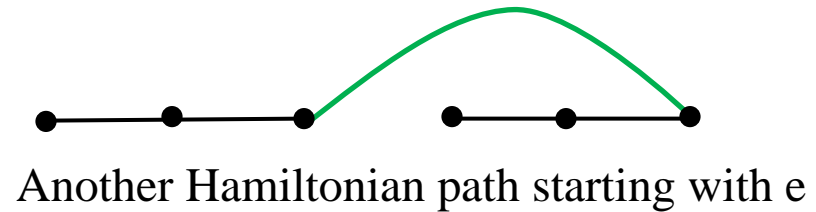
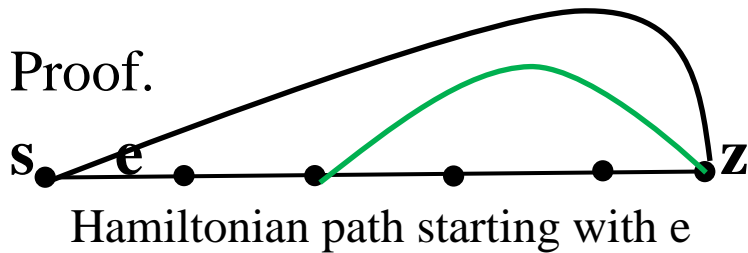
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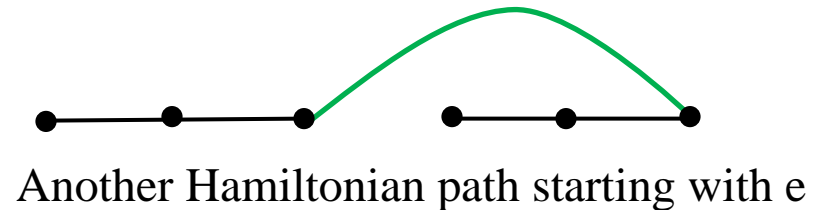
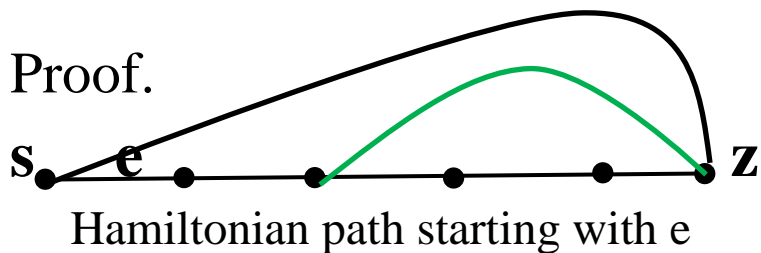
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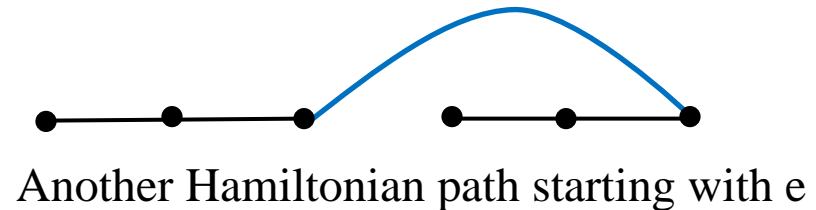
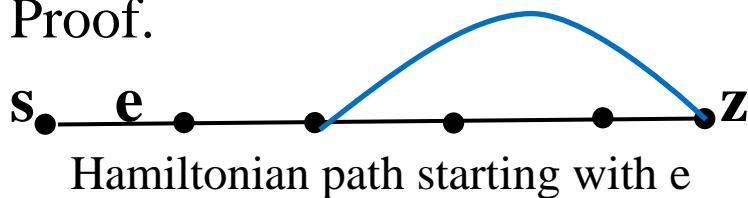
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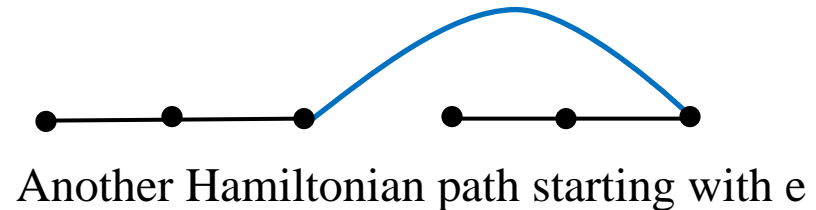
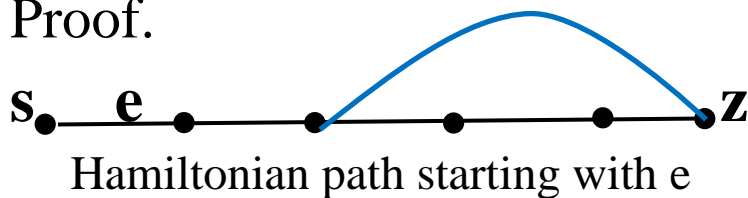
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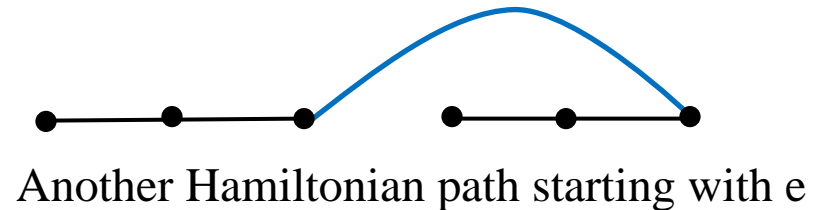
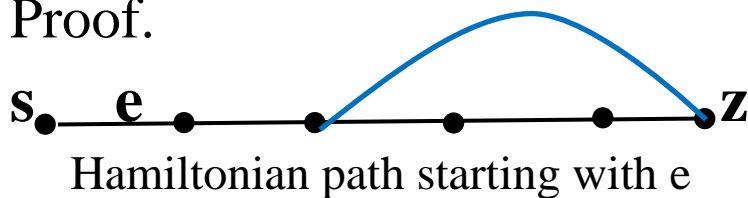
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In any graph, the number of vertices of odd degree is even.

Thus, the number of Hamiltonian cycles containing e is even. \square

Theorem (Shunichi Toida, 1973)

Let G be a graph where the degree, $d(v)$, of every vertex v is even.

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Then the number of cycles containing edge e is odd.

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Since no vertex has odd degree, we could replace “**cycles**” by
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Theorem (Shunichi Toida, 1973)

Let G be a graph where the degree, $d(v)$, of every vertex v is **even**.

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	All vertices have even degree		There is an odd-degree vertex	...	All vertices have odd degree
#cycles containing e and all odd-degree vertices	odd		even	...	even

Carsten Thomassen and I proved that the number of cycles containing e and all the odd-degree vertices is even as soon as the graph has an odd-degree vertex.

A graph is called *Eulerian* if every vertex has even degree.

Theorem (Carsten Thomassen and KC, 2018+)

Let G be a graph and let e be an edge of G .

The number of cycles containing edge e and all the odd-degree vertices is **odd**

if and only if

G is **Eulerian**.

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Let G be a graph and let e be an edge of G .

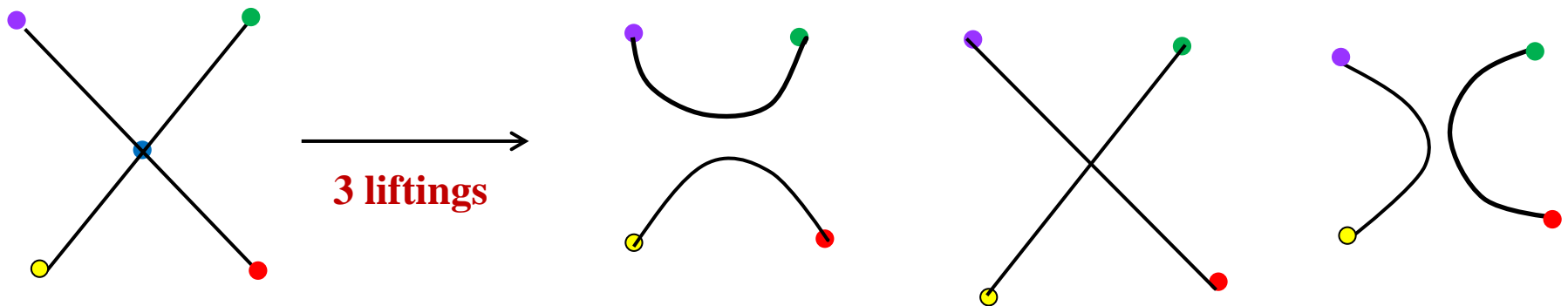
The number of cycles containing edge e and all the odd-degree vertices is **odd**

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Our proof uses “liftings” of even-degree vertices and is not algorithmic.

I will give an exchange graph proof.



Some background:

Theorem (Andrew Thomason 1978)

Let G be a graph where the degree, $d(v)$, of every vertex v is odd. Let e be an edge of G . Then the number of Hamiltonian cycles containing edge e is even.

Corollary

Let G be a graph where the degree, $d(v)$, of every vertex v is odd. Let e be an edge of G . If there is one Hamiltonian cycle containing e , then there is another.

Theorem (Carsten Thomassen, 2016, published 2018)

Let G be a graph where no two even-degree vertices are adjacent. If there is one cycle containing all the odd-degree vertices, then there is another.

Thomassen used four types of exchange operations, one of which was Andrew Thomassen's "lollipop" exchange.

Theorem (KC, 2017+)

Let G be a graph where no two even-degree vertices are adjacent. Let e be an edge of G . Then the number of cycles containing e and all the odd-degree vertices is even.

Follows from

Theorem (KC, 2017+)

Let G be a bipartite graph with bipartition (A,B) , where every vertex in A has odd degree and every vertex in B has even degree. Let e be an edge of G .

Then the number of cycles containing e and all the odd-degree vertices is even.

Some background:

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Theorem (KC, 2017+)

Let G be a graph where no two even-degree vertices are adjacent. Let e be an edge of G . Then the number of cycles containing e and all the odd-degree vertices is even.

But it turned out that by generalizing the exchange operations into two types, one can get an exchange graph proof of a more general result.

Theorem (Carsten Thomassen and KC, 2018+)

Let G be a graph with an odd-degree vertex and let e be an edge of G . The number of cycles containing edge e and all the odd-degree vertices is **even**.

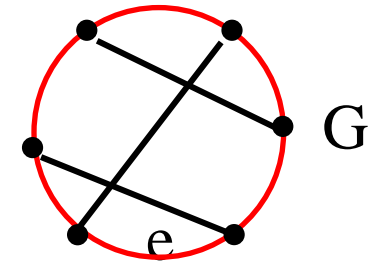
This, combined with Toida's Theorem gives:

Theorem (Carsten Thomassen and KC, 2018+)

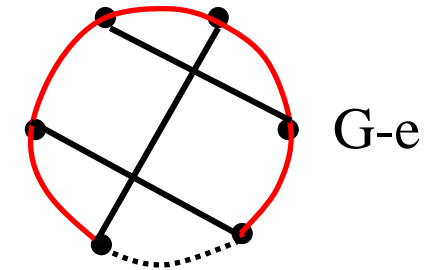
Let G be a graph and let e be an edge of G . The number of cycles containing edge e and all the odd-degree vertices is **odd** if and only if G is **Eulerian**.

Jack Edmonds and I (1999) gave an exchange graph proof of Toida's Theorem.

A Hamiltonian cycle containing edge e in G



corresponds to a Hamiltonian path in $G - e$
which is a spanning tree with degree with degree 1
at the ends of e and 2 elsewhere



Defn. Given a graph G and a subgraph S of G (for us, a tree),
the **excess degree** of a vertex is its degree in $G - E(S)$.

Theorem (Ken Berman, 1986)

Suppose graph G has a spanning tree T where the excess degree of each vertex is odd. Then the number of spanning trees of G with the same degree at each vertex as T is even.

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Suppose graph G has a spanning tree T where the excess degree of each vertex is odd. Then the number of spanning trees of G with the same degree at each vertex as T is even.

Jack Edmonds and I gave an exchange graph proof (1999) of a generalization – only vertices which are not leaves need to have excess degree

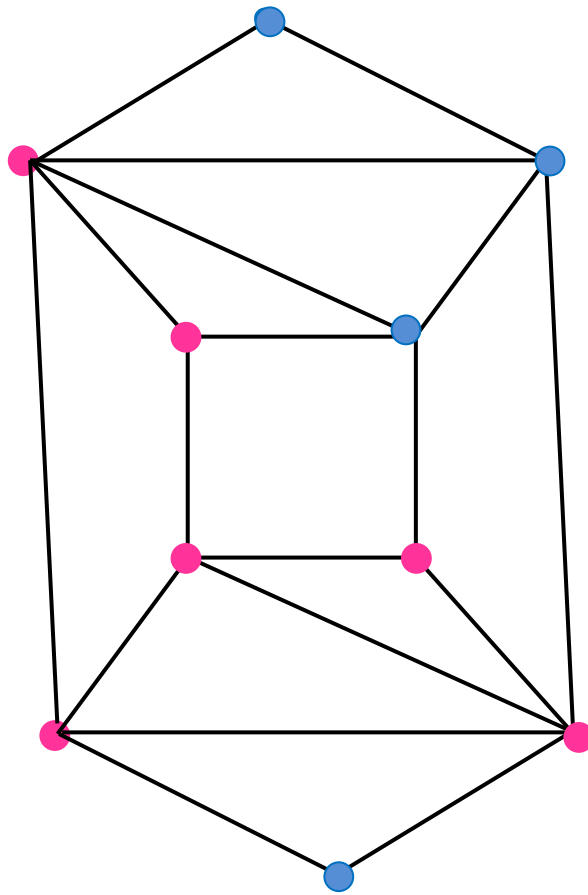
Theorem

*Suppose graph G has a spanning tree T where the excess degree of each vertex **which is not a leaf of T** is odd. Then the number of spanning trees of G with the same degree at each vertex as T is even.*

Let \mathbf{B} be a set of even-degree vertices in a graph G

Let $\mathbf{A} = \mathbf{V}(G) - \mathbf{B}$

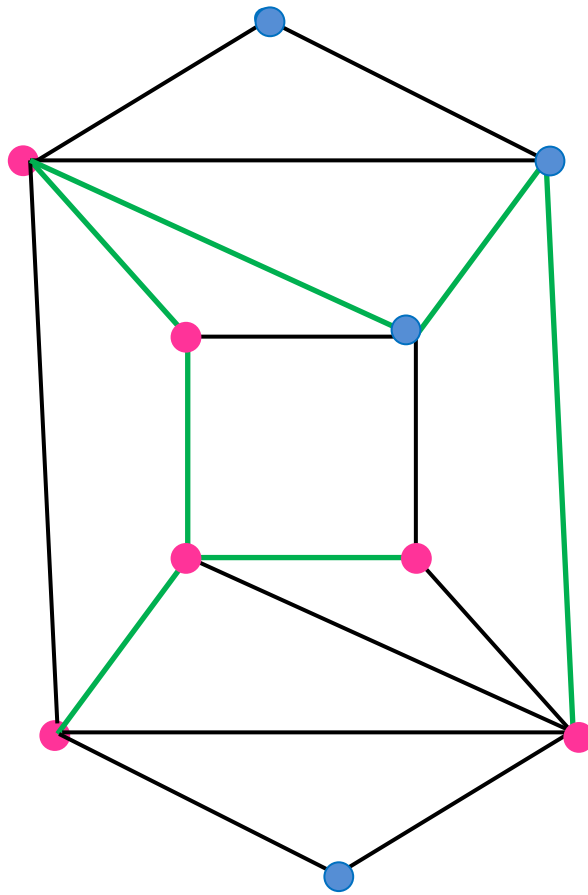
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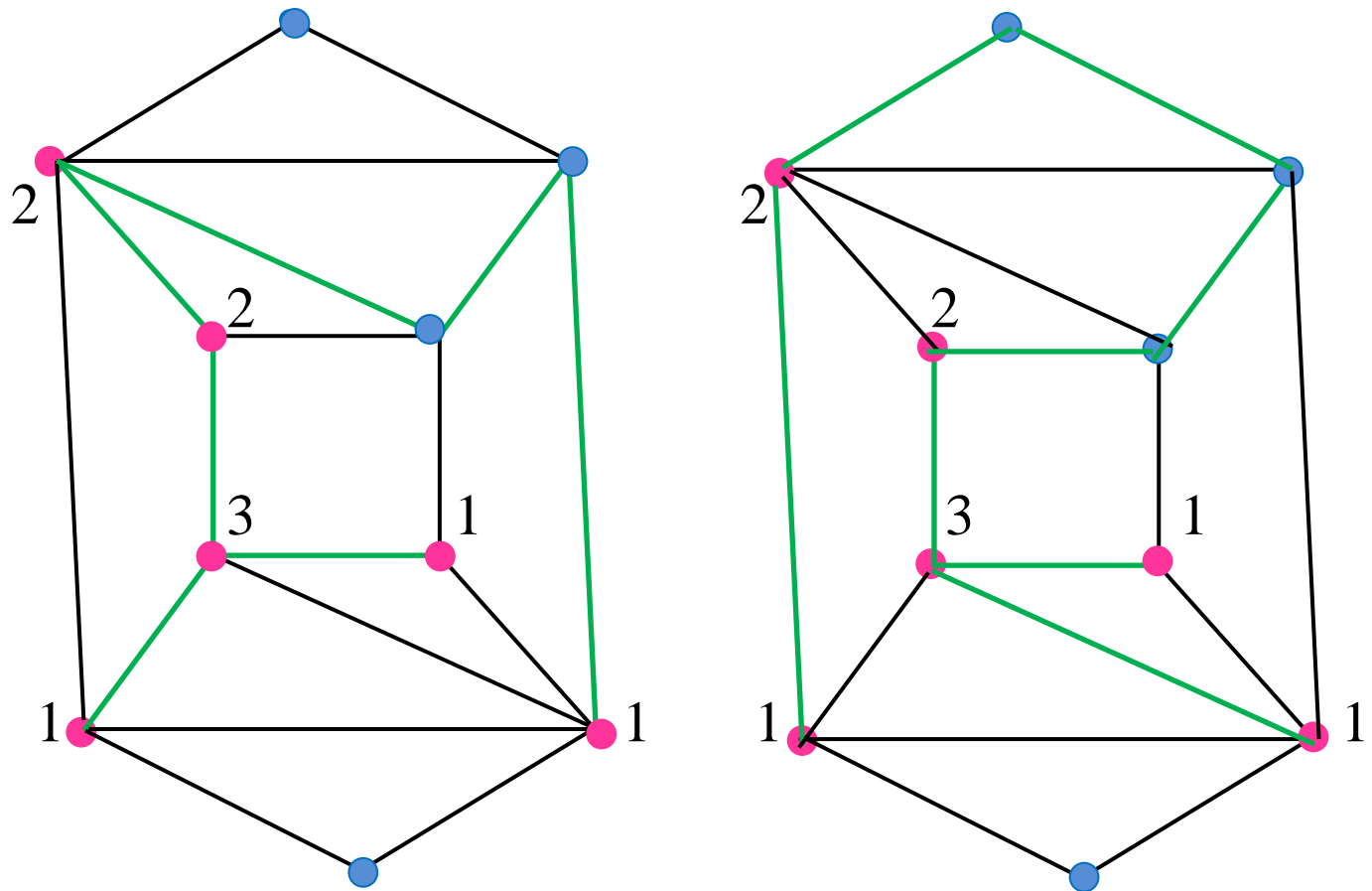


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A good tree T is called \mathbf{T}^* -similar if it has the same degree at each vertex of \mathbf{A} as \mathbf{T}^* .



Let B be a set of even-degree vertices in a graph G . $A = V(G) - B$.

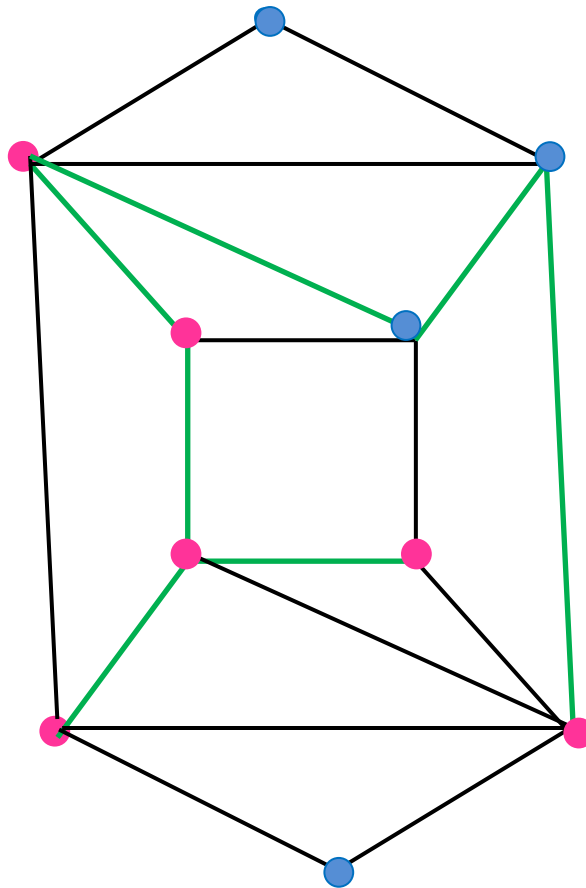
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Consider a good tree T^* such that every vertex of A which is not a leaf

of T^* has odd excess degree



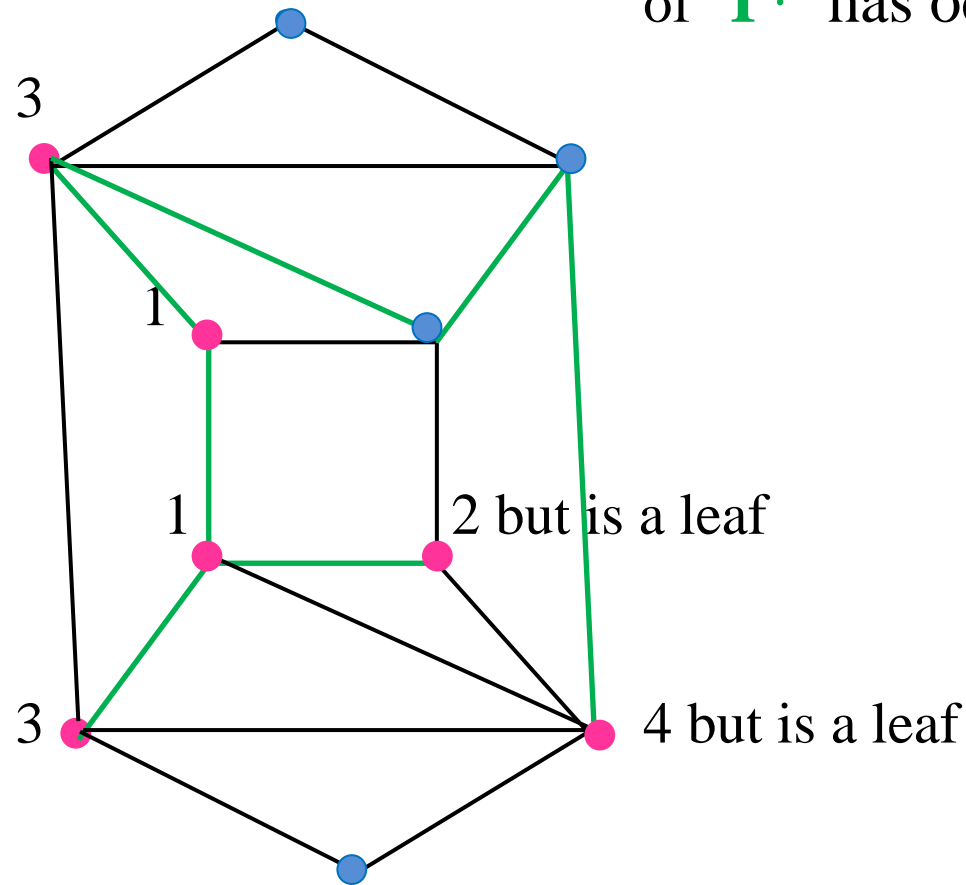
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A good tree T is called T^* -similar if it has the same degree at each vertex of A as T^* .

Theorem (KC 2018+)

Let G be a graph and B a set of even-degree vertices in G .

Let $A = V(G) - B$.

Let T^* be a good tree such that each vertex of T^* which is not a leaf has odd excess degree.

Then the number of T^* -similar trees is even.

Jack Edmonds and I (2017+) had proved this when G is a bipartite graph with bipartition (A, B) .

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Theorem (Carsten Thomassen and KC, 2018+)

Let G be a graph with an odd-degree vertex and let e be an edge of G .

The number of cycles containing edge e and all the odd-degree vertices is **even**.

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Proof of the cycle theorem from the tree theorem.

Let C be a cycle containing edge $e = xy$ and all the odd-degree vertices.

Let B be the set of even-degree vertices in G , other than x and y (which may have odd or even degree). $A = V(G) - B$.

Then $T^* = C - e$ is a tree in a graph $G - e$.

For each vertex of A in T^* other than possibly leaves x and y , its excess degree in $G - e$ is (an odd number $- 2$), thus odd.

By the tree theorem, the number of T^* -similar trees is even.

There is a 1-1 correspondence between these and cycles in G containing e and all the odd-degree vertices.

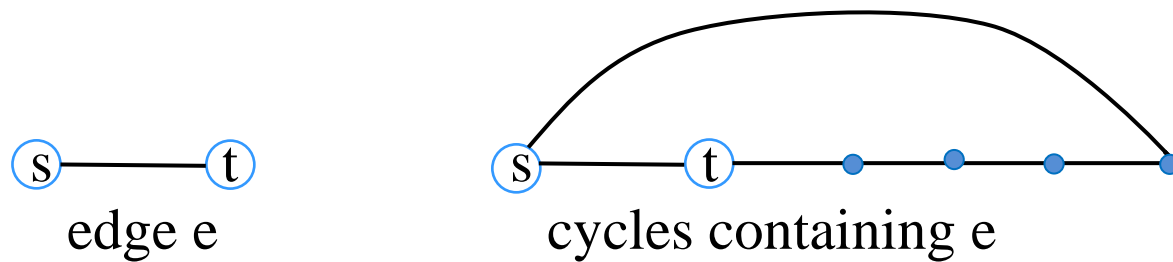
Theorem (Shunichi Toida, 1973)

Let G be a graph where the degree, $d(v)$, of every vertex v is even (i.e. G is Eulerian). Let $e = st$ be an edge of G .

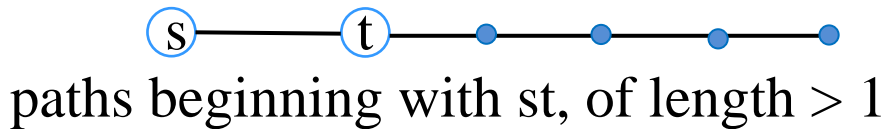
Then the number of cycles containing edge e is odd.

Exchange Graph Proof (Jack Edmonds and KC, 1999)

The odd-degree vertices of the exchange graph $X(G)$ are:



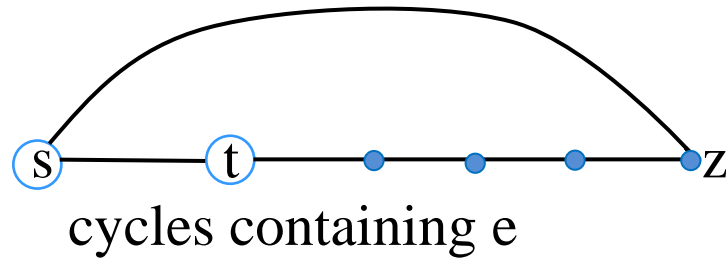
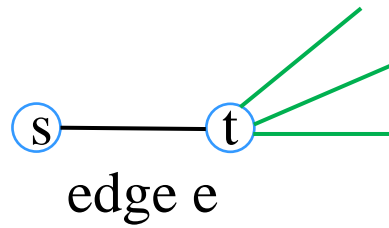
The even-degree vertices of the exchange graph $X(G)$ are:



The exchange operations are adding or removing an edge meeting the last vertex of the path or an Andrew Thomason exchange:



The odd-degree vertices of the exchange graph $X(G)$ are:



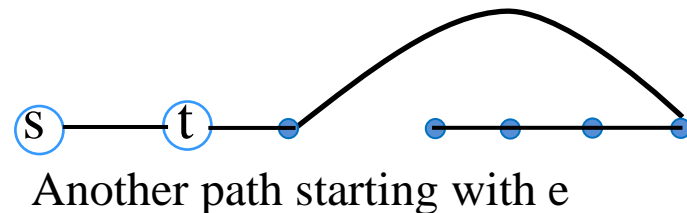
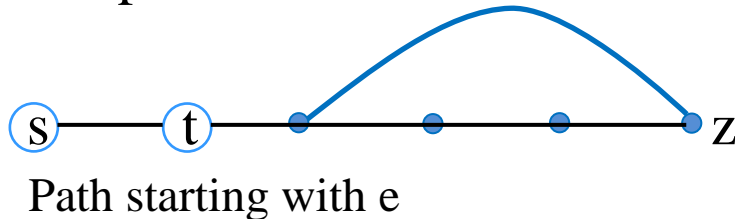
**degree is $d(t)-1$
which is odd**

The even-degree vertices of the exchange graph $X(G)$ are:

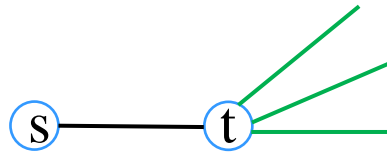


paths beginning with st, of length > 1

The exchange operations are adding or removing an edge meeting the last vertex of the path or an Andrew Thomason exchange:

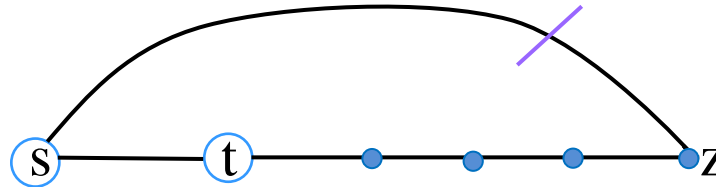


The odd-degree vertices of the exchange graph $X(G)$ are:



edge e

**degree is $d(t)-1$
which is odd**



cycles containing e

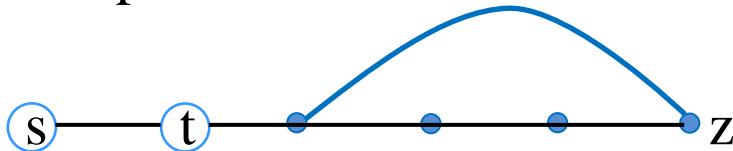
**degree is 1
which is odd**

The even-degree vertices of the exchange graph $X(G)$ are:

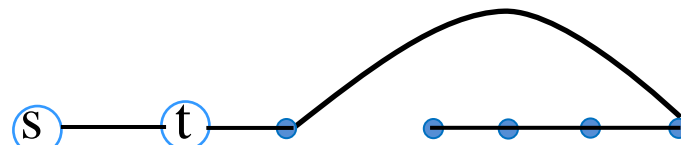


paths beginning with st , of length > 1

The exchange operations are adding or removing an edge meeting the last vertex of the path or an Andrew Thomason exchange:

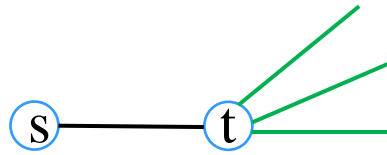


Path starting with e



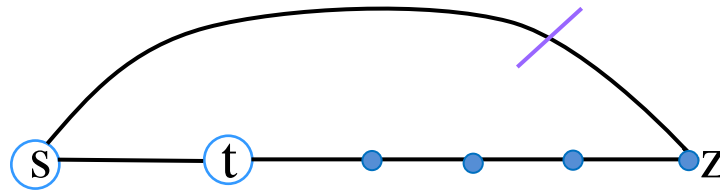
Another path starting with e

The odd-degree vertices of the exchange graph $X(G)$ are:



edge e

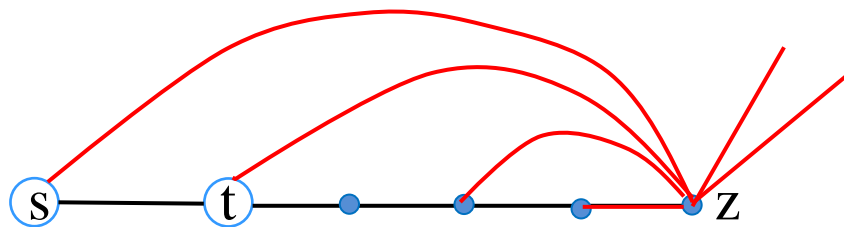
**degree is $d(t)-1$
which is odd**



cycles containing e

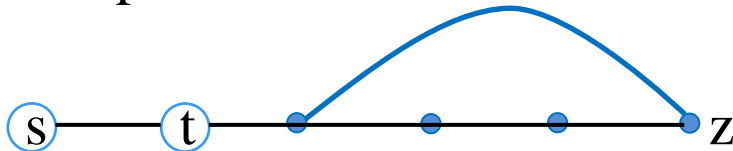
**degree is 1
which is odd**

The even-degree vertices of the exchange graph $X(G)$ are paths beginning with st , of length > 1

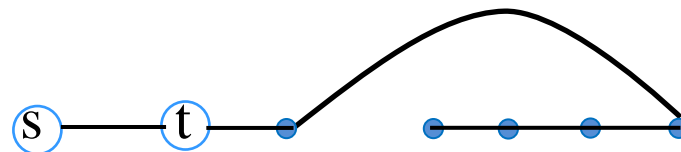


degree is $d(z)$ which is even

The exchange operations are adding or removing an edge meeting the last vertex of the path or an Andrew Thomason exchange:



Path starting with e



Another path starting with e